

Path Optimization and Trusted Routing in MANET: *An Interplay Between Ordered Semirings*

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Motivation

- Trust is a critical objective for routing in MANETs
- Unlike traditional metrics, trust metrics live in abstract ordered sets
 - Consequently, trust metrics cannot be trivially combined with other routing metrics such as delay.
- Multi-criteria optimization techniques are needed for trusted routing

Contributions

- Modeling the Delay-Trust routing problem as a Pareto optimality problem.
 - Combining shortest path and spanning tree solutions
- General optimization principle – in combining (min,sum) and (min,max) algebras
 - Polynomial time algorithms

Outline

- Pareto optimality
- Delay-Trust routing
- Haimes method for Delay-Trust routing

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Vector valued function

 \mathcal{X}

Decision Set

 \mathcal{Y}

Objective Set

 $\underline{f} : \mathcal{X} \rightarrow \mathcal{Y}$

Vector valued map

$$\underline{f} = \begin{bmatrix} f_1 \\ f_2 \end{bmatrix}$$

For the bi-objective case

$$\mathcal{Y} = S1 \times S2$$

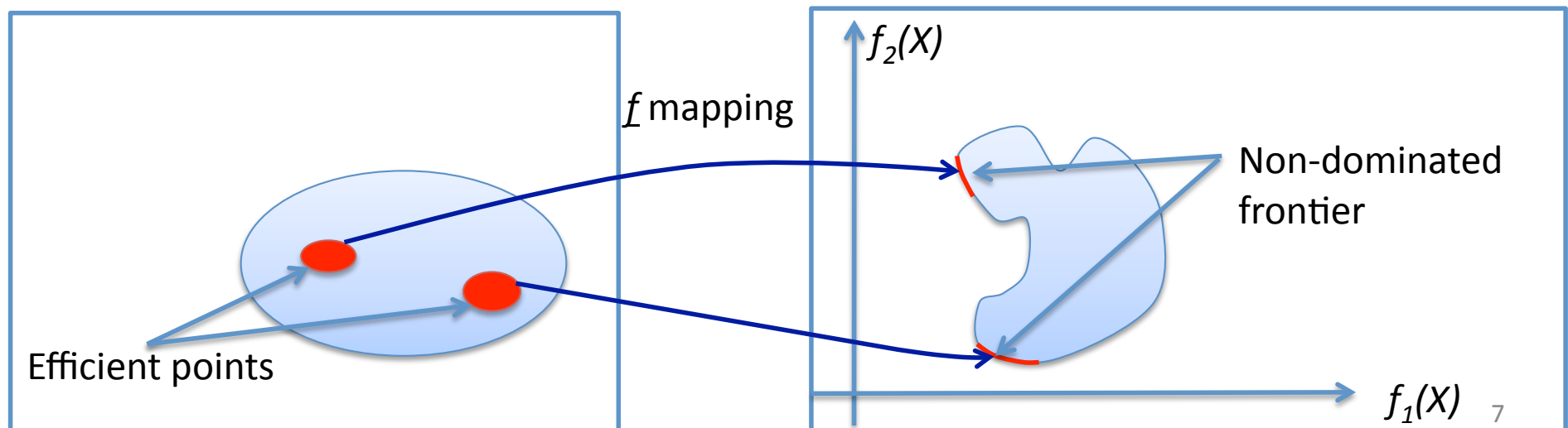
Where S1 and S2 are two ordered sets.

Note: The orders might not be commensurable

Minimality, Non-Dominance, Efficiency

$\underline{f} : X \rightarrow \mathbb{R}^n$ is a vector-valued function.

$x^* \in X$ is efficient if there exists no other $x \in X$ such that $\underline{f}(x) < \underline{f}(x^*)$. $\underline{f}(x^*)$ is said to be non-dominated



ε -relaxed Haimes method

- Relax all but one of the objectives
- With the other objectives as constraints, solve a single objective function.

$$\begin{array}{ll} \min_{x \in X} & f_1(x) \\ \text{subject to} & f_2(x) \leq \varepsilon \end{array}$$

Outline

- Pareto optimality
- **Delay-Trust routing**
- Haimes method for Delay-Trust routing

Pareto Optimal Trusted Routing

- Consider two objectives for routing
 - Delay – additive along a path p

$$f_1(p) = \sum_{(u,v) \in p} d(u,v)$$

- Dual Trust – bottleneck along a path p

$$f_2(p) = \max_{(u,v) \in p} t(u,v)$$

Pareto Optimal Trusted Routing

- Consider two objectives for routing
 - Pareto optimal Delay – Trust problem

$$\begin{bmatrix} f_1(p) \\ f_2(p) \end{bmatrix} = \begin{bmatrix} \sum_{(u,v) \in p} d(u,v) \\ \max_{(u,v) \in p} t(u,v) \end{bmatrix}$$

$$\min_{p \in P_{ij}}^{Pareto} \begin{bmatrix} f_1(p) \\ f_2(p) \end{bmatrix}$$

Pareto Optimal Trusted Routing

- Two objectives for routing – Two classical problems

- Delay – shortest path problem – **(min,+)**

$$\min_{p \in P_{ij}} f_1(p) = \min_{p \in P_{ij}} \sum_{(u,v) \in p} d(u,v)$$

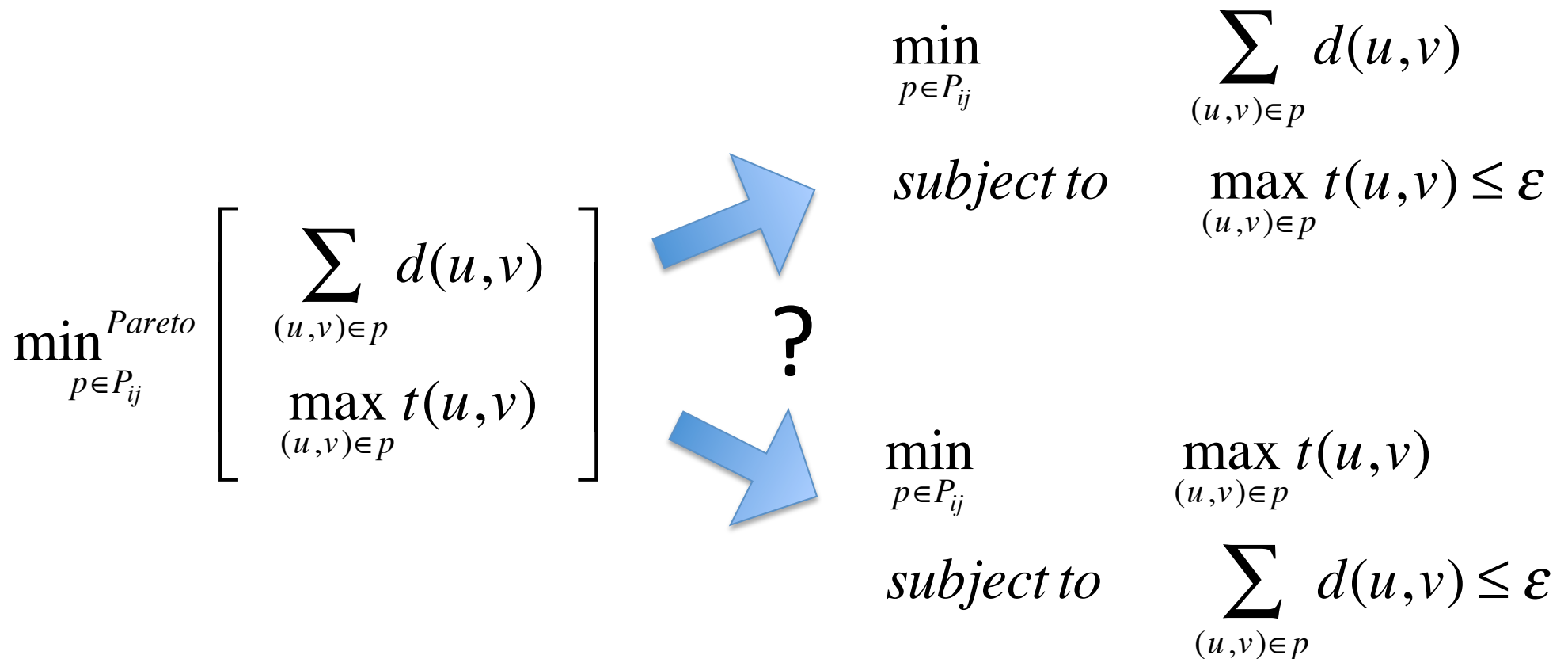
- Dual Trust – spanning tree problem – **(min,max)**

$$\min_{p \in P_{ij}} f_2(p) = \min_{p \in P_{ij}} \max_{(u,v) \in p} t(u,v)$$

Outline

- Pareto optimality
- Delay-Trust routing
- Haimes method for Delay-Trust routing

Hamies method for Trusted Routing



Hamies method for Trusted Routing

$$\min_{p \in P_{ij}}^{\text{Pareto}} \left[\begin{array}{l} \sum_{(u,v) \in p} d(u,v) \\ \max_{(u,v) \in p} t(u,v) \end{array} \right]$$



$$\begin{array}{ll} \min_{p \in P_{ij}} & \sum_{(u,v) \in p} d(u,v) \\ \text{subject to} & \max_{(u,v) \in p} t(u,v) \leq \varepsilon \end{array}$$

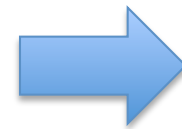
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Hamies method for Trusted Routing

$$\begin{array}{ll} \min_{p \in P_{ij}} & \sum_{(u,v) \in p} d(u,v) \\ \text{subject to} & \max_{(u,v) \in p} t(u,v) \leq \varepsilon \end{array}$$

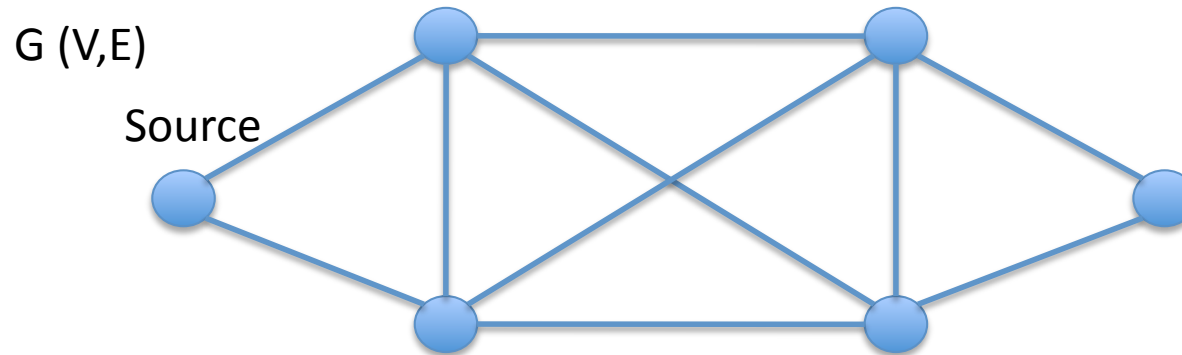
○ Constraint

$$\begin{array}{l} \max_{(u,v) \in p} t(u,v) \leq \varepsilon \\ \Rightarrow \forall (u,v) \in p, \quad t(u,v) \leq \varepsilon \end{array}$$



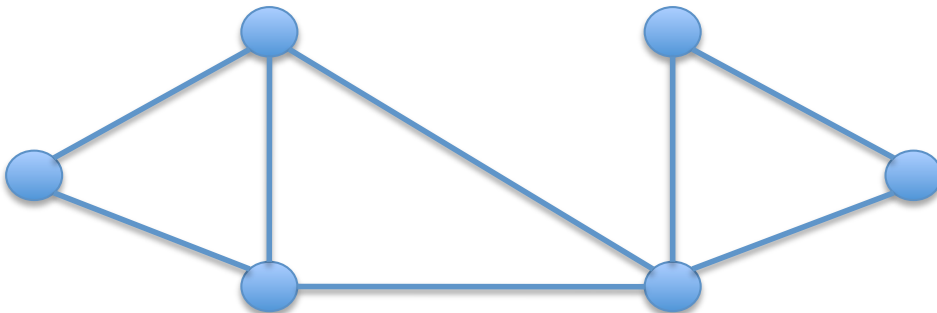
Edge
Exclusion

Haimes method – Two stage recipe



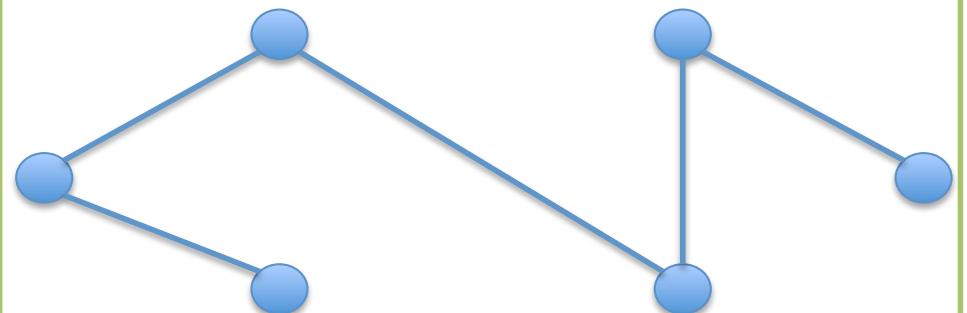
1. G' – reduced graph

$O(|E|)$



2. G'^{SP} – SP on reduced graph

$O(|V| \cdot |E|)$



Summary

- Delay-Trust routing problem can be modeled as a Pareto optimal routing problem.
- Haimes method can be used to solve for Pareto optimality.
 - First solve the (min,max) problem and then the (min,+) problem
 - $O(|V|^3)$ solution